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Nonbinary Spatially-Coupled LDPC Codes on the Binary Erasure Channel

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Abstract—We analyze the asymptotic performance of nonbinary spatially-coupled low-density parity-check (SC-LDPC) codes built on the general linear group, when the transmission takes place over the binary erasure channel. We propose an efficient method to derive an upper bound to the maximum a posteriori probability (MAP) threshold for nonbinary LDPC codes, and observe that the MAP performance of regular LDPC codes improves with the alphabet size. We then consider nonbinary SC-LDPC codes. We show that the same threshold saturation effect experienced by binary SC-LDPC codes occurs for the nonbinary codes, hence we conjecture that the BP threshold for large termination length approaches the MAP threshold of the underlying regular ensemble.

I. INTRODUCTION

Low-density parity-check (LDPC) codes are a powerful class of codes achieving rates very close to capacity for binary memoryless symmetric (BMS) channels. Their excellent performance, however, usually requires long block lengths over the binary field. For short-to-moderate block lengths, nonbinary LDPC codes have been shown to outperform their binary counterparts [1]. For this reason, nonbinary LDPC codes designed over Galois fields of order 2^m (GF_2^m), where m is the number of bits per symbol, have received a considerable interest in the last few years. Their performance under iterative decoding was analyzed in [2], [3]. In [2] the density evolution (DE) for nonbinary LDPC code ensembles defined with respect to the general linear group over the binary field was derived for the binary erasure channel (BEC). It was shown that the messages exchanged in the belief propagation (BP) decoder can be interpreted as subspaces of the vector space GF_2^m , which need to be enumerated. Also, it was observed in [2] that the BP threshold of some nonbinary LDPC code ensembles improves up to a certain m and then worsens for increasing values of m . Upper bounds to the maximum a posteriori probability (MAP) thresholds of nonbinary LDPC code ensembles were also given, and they were conjectured to be tight. Later, in [4], the Maxwell construction of [5], relating the performance of the MAP and the BP decoder, was shown to hold for nonbinary LDPC codes over GF_2^2 .

Spatially-coupled LDPC (SC-LDPC) codes [6] have received notable attention in the recent years due to their outstanding performance for a myriad of channels and communication problems. For the BEC, it was proved in [7] that the BP decoding threshold of a binary SC-LDPC code approaches the optimal MAP threshold of the underlying LDPC block code ensemble, a phenomenon known as *threshold saturation*.

This result has been recently extended to BMS channels [8], and the same phenomenon has been observed for many other channels, such as the multiple access and the relay channel. However, despite their excellent performance for long blocks, for fixed block length SC-LDPC codes perform poorly for the short-to-moderate block length regime, worse than, e.g., irregular LDPC codes.

In this paper, we consider nonbinary SC-LDPC codes for transmission over the BEC. To the best of our knowledge, only the recently submitted paper [9] addresses the construction of nonbinary SC-LDPC codes and reports some BP thresholds for the BEC. However, no analysis on the MAP threshold nor on the threshold saturation phenomenon of nonbinary SC-LDPC codes is performed in [9]. Here, to analyze the asymptotic performance of nonbinary SC-LDPC codes we first consider nonbinary LDPC codes and, in particular, analyze their MAP threshold. We give a systematic and elegant way to generate all subspaces of GF_2^m of a certain dimension, which allows us to compute the BP extrinsic information transfer (BP EXIT) function for an arbitrary m and subsequently an upper bound to the MAP threshold for nonbinary LDPC codes. We show that the MAP threshold bound of regular ensembles improves with m and approaches the Shannon limit. We then analyze nonbinary SC-LDPC codes and show that, contrary to regular and irregular LDPC codes for which the BP decoding threshold worsens for high values of m , the BP threshold of nonbinary SC-LDPC codes with large termination length improves with m (for the values analyzed) and tends to the Shannon limit. We also show the threshold saturation phenomenon for given m , and we conjecture that the BP threshold of SC-LDPC codes with increasing termination length saturates to the MAP threshold of the underlying regular ensemble.

II. NONBINARY LDPC CODES

We consider transmission over a BEC with erasure probability ε using nonbinary LDPC codes defined over the general linear group [2]. The code symbols are elements of the binary vector space GF_2^m of dimension m . The code block length is n symbols, and we transmit on the BEC the m -tuples representing their binary image. Therefore, we interpret the codeword as a binary codeword of length nm . We denote by x_i the i th information bit and by y_i the corresponding channel output, which is a random variable over $\{0, 1, ?\}$, where symbol $?$ denotes an erasure. The channel outputs

are collected into the vector $\mathbf{y} = \{y_1, \dots, y_{nm}\}$. Also, we denote by $\mathbf{y}_{\sim i}$ the vector of the channel outputs when the i th sample is omitted. We denote a regular nonbinary LDPC code ensemble as $\mathcal{G}(d_v, d_c, m)$, where d_v is the variable-node degree and d_c is the check-node degree. Given a code in this ensemble, we associate to each edge of the corresponding graph a bijective linear mapping $f : \text{GF}_2^m \rightarrow \text{GF}_2^m$, chosen uniformly at random. The set of mappings is the general linear group GL_2^m over the binary field, which is the set of all $m \times m$ invertible matrices whose entries take values on $\{0, 1\}$. The design rate r of a code in the ensemble $\mathcal{G}(d_v, d_c, m)$ does not depend on m and can be expressed as $r = 1 - \frac{d_v}{d_c}$.

In this work, we are interested in the asymptotic average performance of the regular ensemble when $n \rightarrow \infty$. The asymptotic performance of LDPC codes can be analyzed in terms of the MAP and BP thresholds. We denote the MAP and the BP thresholds by ε^{MAP} and ε^{BP} , respectively. In the case of transmission over the BEC, ε^{MAP} is the largest channel parameter such that the normalized conditional entropy converges to zero. The evaluation of ε^{MAP} is not an easy task, but an upper bound can be obtained by computing the asymptotic average BP EXIT function, which corresponds to running a BP decoder on a very large graph until the decoder has reached a fixed point. This can be accomplished by means of the DE method [10]. Given a code G in the ensemble \mathcal{G} , the BP EXIT function at the ℓ th iteration is defined as

$$h_G^{\text{BP}, \ell}(\varepsilon) = \frac{1}{nm} \sum_i P(\hat{x}_i^\ell = ? | \mathbf{y}_{\sim i}),$$

where \hat{x}_i^ℓ is the estimate delivered by the BP decoder at the ℓ th iteration. The asymptotic average BP EXIT function of the ensemble is defined as

$$h^{\text{BP}}(\varepsilon) = \lim_{\ell \rightarrow \infty} \lim_{n \rightarrow \infty} \mathbb{E}_G[h_G^{\text{BP}, \ell}(\varepsilon)]. \quad (1)$$

Starting at $\varepsilon = 0$, the function is zero until $\varepsilon = \varepsilon^{\text{BP}}$, at which point it jumps to a non-zero value and continues smoothly until it reaches one at $\varepsilon = 1$. An upper bound for the MAP threshold ε^{MAP} can then be obtained by searching the unique value in $[\varepsilon^{\text{BP}}, 1]$ such that $\int_{\varepsilon^{\text{BP}}}^1 h^{\text{BP}}(\varepsilon) d\varepsilon = r$. Operationally, we integrate the curve $h^{\text{BP}}(\varepsilon)$ starting at $\varepsilon = 1$ until the area under the curve is equal to the design rate of the code. Since we consider regular ensembles whose BP EXIT functions jump at most once, this bound is conjectured to be tight [5]. (For general ensembles, a tighter bound can be achieved by using the extended BP EXIT function [5].)

A. Density evolution of nonbinary LDPC codes

The messages exchanged in the BP decoding are real vectors of length 2^m , the i th element of which represents the a posteriori probability that the symbol is i . In [2], it was shown that in the case of transmission over the BEC the performance does not depend on the transmitted codeword and hence without loss of generality the transmission of the all-zero codeword can be considered. Under this assumption, the messages arising in the BP decoder assume a simplified form. In particular, the non-zero entries of a message are all equal

and the message itself is equivalent to a subspace of GF_2^m . The number of different subspaces of dimension k of GF_2^m is given by the Gaussian binomial coefficient,

$$G_{m,k} = \begin{bmatrix} m \\ k \end{bmatrix} = \begin{cases} 1 & \text{if } k = m \text{ or } k = 0, \\ \prod_{\ell=0}^{k-1} \frac{2^m - 2^\ell}{2^k - 2^\ell} & \text{otherwise.} \end{cases} \quad (2)$$

Since the non-zero elements of a message are equal, it is sufficient to keep track of the dimension of the messages [2]. We say that a message has dimension k if it has 2^k non-zero elements. If a message coming from a node has dimension k , it means that the symbol is known to be one out of 2^k possible symbols or, equivalently, that at that node $m - k$ relations on the bits composing the symbol are known. Let us consider the three subspaces of dimension one of GF_2^2 , $S_1 = \{00, 01\}$, $S_2 = \{00, 10\}$ and $S_3 = \{00, 11\}$. Subspaces S_1 and S_2 are representative of the case where one bit has been recovered and the other is still erased, while S_3 represents the case where the two bits are erased but their sum modulo-2 is known.

Let $P_c^{(\ell)}(k, d_c)$ be the probability that a randomly chosen message computed by a check node and directed to a connected variable node at the ℓ th iteration has dimension k , and let $P_v^{(\ell)}(k, d_v)$ be the probability that a randomly chosen message computed by a variable node and directed to a check node at the ℓ th iteration has dimension k . At the check nodes, the BP decoder computes the sum of the subspaces corresponding to the incoming messages. We have the following recursion for $d_c \geq 3$ and $c = 4, \dots, d_c$ [2]

$$P_c^{(\ell)}(k, 3) = \sum_{i=0}^k \sum_{j=k-i}^k C_{i,j,k}^m P_v^{(\ell)}(i, d_v) P_v^{(\ell)}(j, d_v)$$

$$P_c^{(\ell)}(k, c) = \sum_{i=0}^k \sum_{j=k-i}^k C_{i,j,k}^m P_c^{(\ell)}(i, c-1) P_v^{(\ell)}(j, d_v),$$

where $C_{i,j,k}^m = \frac{G_{m-i, m-k} G_{i, k-j} 2^{(k-i)(k-j)}}{G_{m, m-j}}$ is the probability of choosing a subspace of dimension j whose sum with a subspace of dimension i has dimension k .

At variable nodes, the decoder computes the intersection of the subspaces corresponding to the incoming messages. We denote by $P_\varepsilon(i)$ the probability that the message coming from the channel has dimension i , which is equivalent to the probability that i bits are erased by the channel. We have

$$P_\varepsilon(i) = \binom{m}{i} \varepsilon^i (1 - \varepsilon)^{m-i}, \quad i = 0, \dots, m.$$

We have the following recursion for $d_v \geq 2$ and $v = 3, \dots, d_v$ [2]

$$P_v^{(\ell+1)}(k, 2) = \sum_{i=k}^m \sum_{j=k}^{m-i+k} V_{i,j,k}^m P_\varepsilon(i) P_c^{(\ell)}(j, d_c)$$

$$P_v^{(\ell+1)}(k, v) = \sum_{i=k}^m \sum_{j=k}^{m-i+k} V_{i,j,k}^m P_v^{(\ell+1)}(i, v-1) P_c^{(\ell)}(j, d_c),$$

where $V_{i,j,k}^m = \frac{G_{i,k}G_{m-i,j-k}2^{(i-k)(j-k)}}{G_{m,j}}$ is the probability of choosing a subspace of dimension j whose intersection with a subspace of dimension i has dimension k .

The asymptotic BP threshold is the largest channel parameter such that the decoding is successful and can be found as

$$\varepsilon^{\text{BP}} = \sup\{\varepsilon \in [0, 1] : P_v^{(\ell)}(0, d_v) \xrightarrow{\ell \rightarrow \infty} 1\}.$$

III. BP-EXIT FUNCTION AND MAP THRESHOLD

To draw the asymptotic average BP EXIT function from (1) we need to compute the bit erasure probabilities of the extrinsic BP decoder $P(\hat{x}_i^\ell = ? | \mathbf{y}_{\sim i})$. In this section, starting from the DE equations in the previous section, we obtain the expression of the extrinsic messages delivered by the BP decoder. We then propose a method to compute the extrinsic bit probability $P(\hat{x}_i^\ell = ? | \mathbf{y}_{\sim i})$ to draw the BP EXIT function for arbitrary m .

We can obtain the extrinsic symbol estimate of the BP decoder $\Psi_{\text{ext}}^{(\ell)}$ at the ℓ th iteration taking into account all incoming messages to a variable node from the connected check nodes. Notice that due to the extrinsic nature of the message, the channel observations do not contribute in its computation. We define $P_{\text{ext}}^{(\ell)}(k)$, $k = 0, \dots, m$, the probability that the message $\Psi_{\text{ext}}^{(\ell)}$ has dimension k and we have the following recursion for $d_v \geq 2$ and $v = 3, \dots, d_v$

$$P_{\text{ext}}^{(\ell+1)}(k, 2) = \sum_{i=k}^m \sum_{j=k}^{m-i+k} V_{i,j,k}^m P_c^{(\ell)}(i, d_c) P_c^{(\ell)}(j, d_c)$$

$$P_{\text{ext}}^{(\ell+1)}(k, v) = \sum_{i=k}^m \sum_{j=k}^{m-i+k} V_{i,j,k}^m P_{\text{ext}}^{(\ell+1)}(i, v-1) P_c^{(\ell)}(j, d_c).$$

Finally, we have

$$P_{\text{ext}}^{(\ell+1)}(k) = P_{\text{ext}}^{(\ell+1)}(k, d_v).$$

To compute (1), we let the number of decoder iterations go to infinity. We define the following asymptotic quantities

$$P_{\text{ext}}(k) = \lim_{\ell \rightarrow \infty} P_{\text{ext}}^{(\ell)}(k)$$

$$\hat{x}_i = \lim_{\ell \rightarrow \infty} \hat{x}_i^{(\ell)}, \quad \Psi_{\text{ext}} = \lim_{\ell \rightarrow \infty} \Psi_{\text{ext}}^{(\ell)}$$

and obtain

$$P(\hat{x}_i = ? | \mathbf{y}_{\sim i}) = \sum_{k=0}^m P(\hat{x}_i = ? | \dim(\Psi_{\text{ext}}) = k, \mathbf{y}_{\sim i}) P_{\text{ext}}(k). \quad (3)$$

Thanks to the tree assumptions, the probabilities in (3) do not depend on i and (3) is actually the BP extrinsic entropy of a bit $h^{\text{BP}}(\varepsilon)$. To evaluate the probabilities $P(\hat{x}_i = ? | \dim(\Psi_{\text{ext}}) = k, \mathbf{y}_{\sim i})$, we have to enumerate the subspaces associated to the message Ψ_{ext} , so that we can compute

$$P(\hat{x}_i = ? | \dim(\Psi_{\text{ext}}) = k, \mathbf{y}_{\sim i}) = \frac{1}{G_{m,k}} \sum_{z=1}^{G_{m,k}} P_z^{(k)}(\hat{x}_i = ? | \mathbf{y}_{\sim i}),$$

where $P_z^{(k)}(\hat{x}_i = ? | \mathbf{y}_{\sim i})$ is the extrinsic bit erasure probability associated to the z th subspace of dimension k , given an arbitrary but fixed ordering.

We propose an efficient method to identify the subspaces and to derive the corresponding extrinsic erasure probability. Note that each subspace of dimension k can be interpreted as the set of 2^k codewords of an (m, k) binary linear block code of length m . Therefore, we can associate to each subspace of dimension k an $(m-k) \times m$ matrix (the parity-check matrix of the code) such that the symbols corresponding to the non-zero entries in the message belong to the nullspace of the matrix, meaning that the subspace associated to the message is the nullspace of the matrix. Thus, finding all subspaces of a certain dimension k reduces to find all parity-check matrices that generate a different (m, k) code. An efficient way to find these matrices is to find all $(m-k) \times m$ binary matrices in row-reduced echelon form containing no zero rows. The set of these matrices is denoted here as $\mathcal{R}_{m,k}$. A row-reduced echelon binary matrix is defined as a matrix in which (i) the first one in every row is in a column where all other elements are zero and (ii) the number of leading zeros increases in every row. The nature of these matrices ensures that their nullspaces are distinct. Thus, the number of matrices in $\mathcal{R}_{m,k}$ corresponds to the number of different subspaces (or codes) of dimension k of GF_2^m . These matrices can be efficiently found by using a modified version of the algorithm described in [11].

Example 1: There are seven different subspaces of dimension one of GF_2^3 , and we associate to them the matrices in $\mathcal{R}_{3,1}$

$$\mathcal{R}_{3,1} = \left\{ \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}, \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \right.$$

$$\left. \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}, \begin{bmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \right\}.$$

The next step consists of the computation of the erasure probability associated to each subspace. We know that $P(\hat{x}_i = ? | \dim(\Psi_{\text{ext}}) = 0, \mathbf{y}_{\sim i}) = 0$ since if the subspace has dimension 0 the symbol is perfectly recovered after BP decoding. Also, $P(\hat{x}_i = ? | \dim(\Psi_{\text{ext}}) = m, \mathbf{y}_{\sim i}) = 1$ since in this case we have complete uncertainty on the transmitted symbol. The cases $k = 1, \dots, m-1$ are less intuitive. However, this problem is easy to solve by interpreting each subspace as an (m, k) code with corresponding parity-check matrix in $\mathcal{R}_{m,k}$ and computing the probability of erasure assuming transmission over the BEC with erasure probability ε . Since the length of these codes is generally very short, we can evaluate the erasure probability by using the full complexity decoding algorithm. Given the z th subspace of dimension k , whose nullspace is generated by the matrix R^z we compute $P_z^{(k)}(\hat{x}_i = ? | \mathbf{y}_{\sim i})$ as [10]

$$\sum_{\mathcal{E} \subseteq [m] \setminus \{i\}} \varepsilon^{|\mathcal{E}|} (1 - \varepsilon)^{m-1-|\mathcal{E}|} (1 + \text{rank}(R_{\mathcal{E}}^z) - \text{rank}(R_{\mathcal{E} \cup i}^z)),$$

where $\mathcal{E} \subseteq [m] = \{1, \dots, m\}$ denotes the index set of erasures and $R_{\mathcal{E}}^z$ denotes the submatrix of R^z indexed by the elements

TABLE I
EXPRESSION OF $h^{\text{BP}}(\varepsilon)$ FOR $m = 1, \dots, 4$. $G_{m,k}$ IS THE GAUSSIAN BINOMIAL COEFFICIENT.

| m | $h^{\text{BP}}(\varepsilon)$ |
|-----|--|
| 1 | $P_{\text{ext}}(1)$ |
| 2 | $(1 + \varepsilon) \frac{P_{\text{ext}}(1)}{G_{1,2}} + P_{\text{ext}}(2)$ |
| 3 | $(1 + 2\varepsilon + \varepsilon^2) \frac{P_{\text{ext}}(1)}{G_{3,1}} + (3 + 4\varepsilon - \varepsilon^2) \frac{P_{\text{ext}}(2)}{G_{3,2}} + P_{\text{ext}}(3)$ |
| 4 | $(1 + 3\varepsilon + 3\varepsilon^2 + \varepsilon^3) \frac{P_{\text{ext}}(1)}{G_{4,1}} + (7 + 18\varepsilon + 9\varepsilon^2 - 6\varepsilon^3) \frac{P_{\text{ext}}(2)}{G_{4,2}} + (7 + 12\varepsilon - 6\varepsilon^2 + \varepsilon^3) \frac{P_{\text{ext}}(3)}{G_{4,3}}$ |

TABLE II
BP AND MAP THRESHOLDS FOR ENSEMBLES $\mathcal{G}(3, 6, m)$.

| m | ε_{BP} | $\bar{\varepsilon}_{\text{MAP}}$ |
|-----|---------------------------|----------------------------------|
| 1 | 0.42944 | 0.48815 |
| 2 | 0.42347 | 0.49487 |
| 3 | 0.41220 | 0.49791 |
| 4 | 0.39890 | 0.49920 |
| 5 | 0.38547 | 0.49970 |
| 6 | 0.37288 | 0.499895 |
| 7 | 0.36154 | 0.499965 |

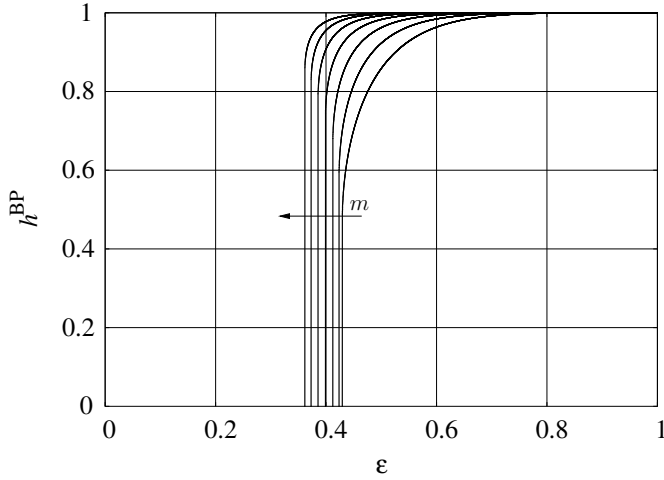


Fig. 1. BP EXIT curves for regular ensembles $\mathcal{G}(3, 6, m)$, with $m = 1, \dots, 7$.

of \mathcal{E} . In Table I, we report the expression of $h^{\text{BP}}(\varepsilon)$ for $m = 1, \dots, 4$.

A. Results

In Fig. 1, we plot the asymptotic average BP EXIT curves for the regular ensembles $\mathcal{G}(3, 6, m)$, with $m = 1, \dots, 7$. We recall that the BP threshold ε_{BP} is the point at which the curve jumps to a non-zero value. The figure shows that the BP threshold decreases as the alphabet size increases. This is in line with previous results in the literature, which state that the performance of the BP decoder degrades by moving to nonbinary alphabets when the variable node degree distribution is greater than two [2]. On the other hand, the MAP performance of the considered ensembles seems to improve when m increases. In Table II, we list the BP thresholds and the upper bounds on the MAP threshold $\bar{\varepsilon}_{\text{MAP}}$. Note that $\bar{\varepsilon}_{\text{MAP}}$ rapidly increases with m and approaches the channel capacity. Similar results were obtained for other ensembles.

IV. NONBINARY SPATIALLY-COUPLED LDPC CODES

We consider nonbinary SC-LDPC code ensembles similar to the ensembles defined in [12] for binary codes, which are derived from regular convolutional protographs by termination.

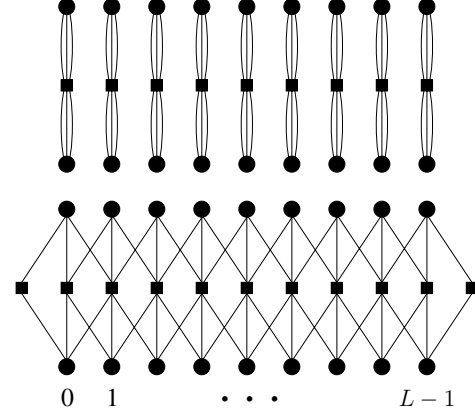


Fig. 2. Chains of nine protographs for the regular ensemble $\mathcal{G}(3, 6, m)$. Top: non-interacting protographs. Bottom: a coupled chain of protographs.

Consider as an example the coupling of regular codes with $d_v = 3$ and $d_c = 6$. The protograph for a $\mathcal{G}(3, 6, m)$ ensemble is composed of two variable nodes and one check node, as shown in Fig. 2 at the top. We consider a chain of L protographs, adding a *spatial* dimension to our code. Since these graphs do not interact, the chain behaves like the original $(3, 6)$ regular code. An SC-LDPC code is then obtained by coupling the protographs: We connect the protograph at position i , $i = 0, \dots, L - 1$, to one neighboring protograph on the left, and to one neighboring protograph on the right, as shown in Fig. 2 at the bottom. We denote this coupled ensemble as $\mathcal{G}_C(3, 6, m, L)$. Locally, the connectivity does not change with respect to the underlying ensemble. The only difference is at the boundaries, where one check node is added on each side to terminate the chain. Check nodes at positions $i \in [1, L - 2]$ have degree six, while the degree of the remaining check nodes decreases linearly according to their position. Having lower degree check nodes helps the decoder, at the expense of a loss in terms of design rate, which is reduced to $r_{3,6}(L) = \frac{1}{2} - \frac{1}{L}$ [7]. By increasing L , the rate loss is reduced, while the beneficial effect of low-degree check nodes on the BP performance does not vanish.

To draw the BP EXIT function for the coupled ensemble, we apply the DE described in Section II-A to each section $i \in [0, L - 1]$, taking into account of the *spatial* structure of the code. We then compute the extrinsic symbol estimate for each section and the corresponding BP extrinsic bit entropy. The BP EXIT function of the SC ensemble is finally obtained by averaging over the L entropies of the chain.

In Fig. 3, we show the BP EXIT curves of the ensembles $\mathcal{G}_C(3, 6, m, L)$ for $m = 1$ and $m = 3$, for several values of L . The corresponding BP and MAP thresholds for $m = 3$ are

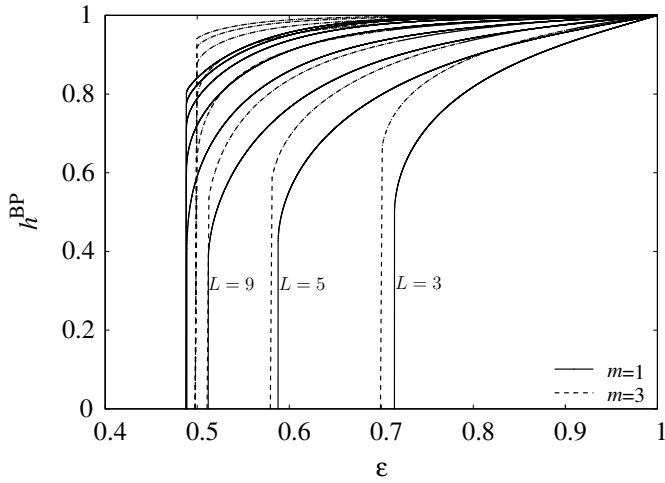


Fig. 3. BP EXIT curves for SC ensembles $\mathcal{G}_C(3, 6, m, L)$. Comparison between binary ($m = 1$) and nonbinary ($m = 3$) codes, for $L = 3, 5, 9, 17, 33, 65, 129$ and 257 .

TABLE III

ASYMPTOTIC THRESHOLDS FOR ENSEMBLES $\mathcal{G}_C(3, 6, 3, L)$ IN FIG. 3.

| L | ε_{BP} | $\tilde{\varepsilon}_{MAP}$ | ε_{Sh} |
|-----|--------------------|-----------------------------|--------------------|
| 3 | 0.69913 | 0.82738 | 0.83333 |
| 5 | 0.57947 | 0.68328 | 0.7 |
| 9 | 0.51077 | 0.59026 | 0.61111 |
| 17 | 0.49795 | 0.54169 | 0.55882 |
| 33 | 0.49791 | 0.51847 | 0.53030 |
| 65 | 0.49791 | 0.50813 | 0.51538 |
| 129 | 0.49791 | 0.50272 | 0.50775 |
| 257 | 0.49791 | 0.50065 | 0.50389 |

given in Table III, where we also report the Shannon limit $\varepsilon_{Sh} = 1 - r_{3,6}(L)$. These results show that the same *threshold saturation* effect observed for binary SC-LDPC codes occurs for the nonbinary codes. In fact, the BP threshold for large L approaches the MAP threshold of the regular ensemble $\mathcal{G}(3, 6, 3)$. Furthermore, since the MAP performance of the underlying ensemble for $m = 3$ outperforms that of the binary one (see Table II), the BP thresholds of the nonbinary SC-LDPC code *saturate* to a better value.

In Fig. 4, we report the BP EXIT curves of the ensembles $\mathcal{G}_C(3, 6, m, 257)$, for $m = 1, \dots, 7$. The figure shows the interesting result that, contrary to non-coupled ensembles, the performance of SC-LDPC codes under BP decoding improves with m and approaches the Shannon limit. Although not reported here due to lack of space, we obtained similar results for other ensembles.

V. CONCLUSIONS

We considered nonbinary SC-LDPC codes over GF_2^m on the BEC. We proposed an efficient method to compute the BP EXIT function for nonbinary LDPC codes and arbitrary alphabet size and subsequently an upper bound to the MAP threshold. Our analysis showed that performance under MAP decoding improves as the alphabet size increases. Furthermore, we analyzed nonbinary SC-LDPC code ensembles, showing that the same threshold saturation effect observed for binary SC-LDPC codes occurs for nonbinary codes. Interestingly, the performance of nonbinary SC-LDPC codes under BP decoding improves with m and approaches the Shannon limit, contrary

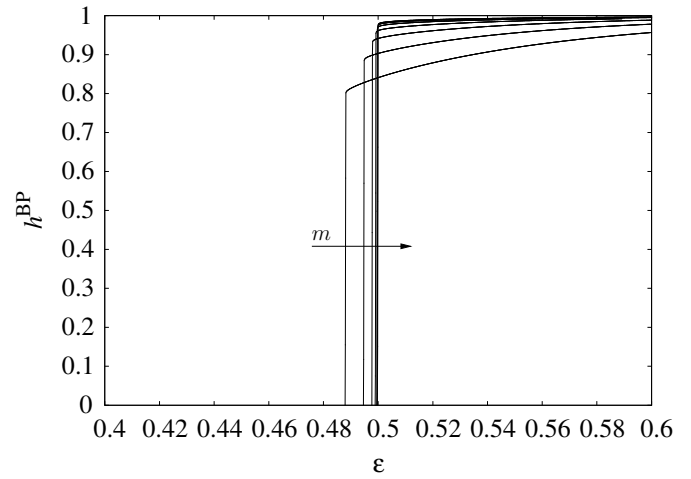


Fig. 4. BP EXIT curves for SC ensembles $\mathcal{G}_C(3, 6, m, 257)$, for $m = 1, \dots, 7$.

to the case of nonbinary, non-coupled LDPC codes. For instance, for the $\mathcal{G}_C(3, 6, m, L)$ ensemble, the BP threshold is improved from 0.4881 for the binary case to 0.4997 for $m = 5$ when L tends to infinity. Future work includes a formal proof of the threshold saturation phenomenon for nonbinary SC-LDPC codes, and a finite-length analysis.

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